Servers in the Mist Operating Systems for Server Utilities

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The Server Cloud



Sharing the Server Cloud



E.g., CNN on 9/11 Facing a World Crisis [William LeFebvre]

The world is dynamic, dangerous, unpredictable, and expensive.

- changing load/traffic
- changing demands
- unexpected failures
- unexpected load surges

Fluid mapping

Load multiplexing Resource efficiency Surge protection Robustness Incremental growth Economy of scale



Outline

Cluster-on-Demand (COD)

- Decentralized resource management for mixed-use clusters
- Automated configuration of dynamic virtual clusters
- Server provisioning for Web server utilities
 - [SOSP-01, USITS-03]
 - Dynamic thermal management (w/ HP Labs)
- Opus (Overlay Peer Utility Service) and ModelNet
 - Self-organizing network of server PoPs [OpenArch-02]
 - Evaluating adaptive Internet systems by emulation [OSDI -02]



Life in Our Shared Cluster



Internet Systems and Storage Group

- Internet software prototypes
- Synthetic load generators
- ModelNet emulations
- Hacked kernels with bugs

Computation Loads

- BioGeometry
- Full-system simulation
- Neuro-engineering simulations
- NC Biogrid PoP (more coming)
- Visualization

Heterogeneous servers Diverse configurations Manual reservation protocol Bursty and deadline-driven

Server Energy



866 MHz P-III SuperMicro 370-DER (FreeBSD), Brand Electronics 21-1850 power meter

Beyond Beowulf

 Compute clusters used to be hard to manage, but we've seen lots of progress since NOW.

- Industry manageability initiatives (PXE, ACPI)
- Configure/monitor/install: NPACI Rocks, Millennium/Rootstock.
- Batch scheduling: LSF, PBS/Maui, SGE, etc.
- Resource sharing/scavenging: Condor, Grids.
- Most assume a homogeneous software base and create a uniform cluster view for batch computation.
 - Beowulf model: Linux/Unix with middleware
- How to extend automated configuration and policy-based scheduling to mixed-use cluster utilities?
 - OS-agnostic? Middleware-agnostic?

COD: Dynamic Virtual Clusters



Enabling Technologies

Linux driver modules, Red Hat *Kudzu,* partition handling

PXE network boot

I P-enabled power units

DHCP host configuration

DNS, NIS, etc. user/group/mount configuration

> NFS etc. network storage automounter

Power: APM, ACPI, Wake-on-LAN

Ethernet VLANs



COD Priming Times

- Time to install images on one node
 - I mage installation will dominate configuration time
- Node originally in "waiting" state
- Bottleneck: disk write speed

Image	Size (KB)	Size (gz)	% Full	Time (s)
CS Linux	2,048,256	105,456	23.2	122
CS BSD	2,048,256	137,020	27.4	129
Debian Linux	1,951,866	422,268	40.6	127

COD and **Related** Systems



COD addresses hierarchical dynamic resource management in mixed-use clusters with pluggable middleware ("multigrid").

A Note on Layering and Hierarchy

Virtual Cluster Manager(VCM middleware)

OS

kernel

COD manager

Virtual machine VMM

process

Like a VMM, COD manages resources outside/beneath the OS and without its knowledge or assistance: "underware".



- Constraint-aware node assignment/negotiation
- Leased allotments for static clusters
- Load-aware provisioning for dynamic clusters
 - Balance local autonomy with global coordination
- Utility functions reflect value and priority (SLAs)
 - "Common currency" weight for VCM demand bids

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Muse: Adaptive Provisioning for a Web Service Utility [SOSP-01]



Multiplex services on nodes, based on proportional-share scheduling in node OS.

Filtering Load Estimates

• EWMA-based filter alone is not sufficient.

- Average A_t for each interval t: $A_t = \alpha A_{t-1} + (1-\alpha)O_t$
- The gain α may be variable or *flip-flop*.
- Load estimate $E_t = E_{t-1} \underline{if} |E_{t-1} A_t| < tolerance$

<u>else</u> $E_t = A_t$

Stable





Cost vs. quality: a simple example Active set Active set **{A,B}** {A, B, C, D} $\rho_i = \rho_{target}$ $\rho_i < \rho_{target}$ • Meets quality goals Low latency • Efficient use

ptarget = configurable target utilization (leaving headroom for transient load spikes).

IBM trace (before)



IBM trace (after)



Maximizing "Revenue" Under Constraint

Input: the "value" of performance for each customer i.

- Common unit of value: "money".
- Derives from the economic value of the service.
- Enables SLAs to represent flexible quality vs. cost tradeoffs.
- Per-customer utility function $U_i = bid + penalty$.
 - Bid for traffic volume (throughput λ_i).
 - Bid for better quality, or add negative *penalty* for poor quality.
- Allocate resources to maximize expected global utility ("revenue" or reward).
 - Model predicts performance effects.
 - "Sell" μ to the highest bidder.
 - Never sell resources below cost.

Maximize $\sum U_i(\lambda_i(t, \mu_i))$ Subject to $\sum \mu_i < \mu_{max}$

SLAs as Utility Functions

Customer SLAs are specified as utility functions

- How much will customer "pay" for a given level of performance (or performability, data quality)?
- Allocate resources to the highest predicted marginal benefit at current load and system conditions.



Resource allotment μ_i

Dynamic Thermal Management



Figure 5: Temperature contour plot at 1.2m above floor: uniform workload, failed CRAC

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Coordinated control of a network of server sites (PoPs) Monitor and adapt for target performance and availability. Configure/instantiate services on demand at selected sites. Connect sites (e.g, replicas) with per-service overlay topology - Metrics: bandwidth, latency, loss rate, cost (\$) sensitivity Decentralized management + weak global coordination

Evaluating Internet Systems

Adaptive Internet software

- Scalable, self-organizing, fast, robust,...
- How to evaluate prototypes?
 - How will it behave under dynamic conditions?
 - How will it behave when the unexpected occurs?
- Simulation
 - There is no substitute for build-and-measure.
- Live deployment
 - Runs real code under real (uncontrolled) conditions: dangerous and expensive, does not yield reproducible results.

The ModelNet Vision

• Emulation environment for adaptive, large-scale services

- Web services, multimedia distribution, p2p, mobile systems
- Accelerate development of robust, adaptive Internet software

• <u>Challenges</u>

- Scalable to 10k nodes, 10 Gb/s bisection bandwidth
 - On 100-node cluster
 - Unmodified applications run on unmodified OS's
- Accurately emulate wide-area network conditions
 - Failures, cross traffic, rapid changes, congestion
- Impossible to capture full complexity of Internet in a machine room
 - Application-specific accuracy vs. scalability tradeoffs



Self-Organizing Overlay



Server Utility in a Dynamic World A Few Key Challenges



- Multiple goals and dimensions of service quality (SLAs)
- Coordinated provisioning of multiple resources that interact in a complex way
- Configure service spread and location for performability
- Reconcile local autonomy and global control

http://issg.cs.duke.edu

http://www.cs.duke.edu/~chase